

An Efficient Message Dissemination Technique in Platooning Applications



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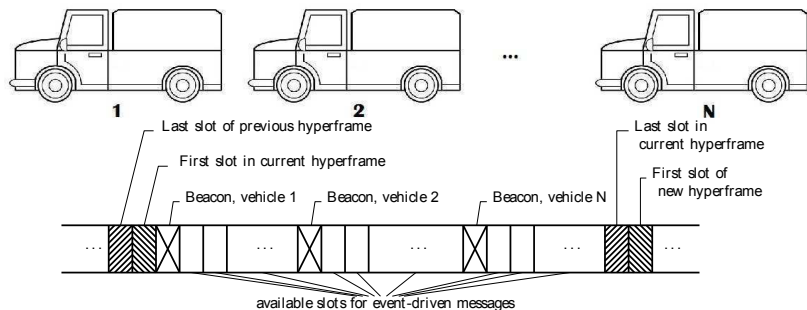
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INTRODUCTION

- ▶ Objective:
 - ▶ *Delay-sensitive* data traffics are sent with high *reliability*
 - ▶ Considering the *co-existence* of both time-triggered and event-driven control messages
- ▶ Current approach:
 - ▶ CSMA IEEE 802.11p: unbounded channel access delay and multiple consecutive packet drops
 - ▶ Time-triggered messages: congestion control, rate control
 - ▶ Event-triggered messages: algorithms to mitigate broadcast storms
 - ▶ Collision-free phase together with contention-based phase, using a separate SCH

SYSTEM MODEL



- ▶ A platoon consisting of N nodes communicating through an *ad hoc* network, using a separate SCH
- ▶ Time is slotted (TDMA) and every packet is transmitted in exactly one time-slot
- ▶ In each superframe, a subset of the available time-slots are dedicated to periodic time-triggered messages

SYSTEM MODEL (CONT'D)

- ▶ Channels between nodes are characterized by an $N \times N$ stationary matrix \mathbf{P}
- ▶ Each event-driven message is assigned a set of $K + 1$ time-slots
- ▶ Time-slot 0 is always allocated to the source R_0 and the remaining K time-slots to a set of K *distinct* relayers R_1, R_2, \dots, R_K
- ▶ After overhearing in time-slots numbered $0, 1, \dots, k - 1$, relay R_k will transmit at time-slot k or being quiet
- ▶ How to choose an optimal combination of R_k , $i = 1, 2, \dots, K$?

ANALYTICAL EXPRESSIONS

- ▶ Defining

$$e_i = \begin{cases} T & \text{if } R_i \text{ transmits at time-slot } i \\ Q & \text{if } R_i \text{ remains quiet at time-slot } i, \end{cases}$$

- ▶ For a specific sequence $\mathbf{e} = \{e_i\}_{i=0}^K \in \{T, Q\}^K$, defining

$\Phi_T(\mathbf{e}) = \{i | e_i = T\}$: the set of time-slots
where transmissions occur

$\Theta_Q(\mathbf{e}) = \{j | e_j = Q\}$: the complimentary set of time slots
where the corresponding relay nodes remain quiet

- ▶ The probability that the destination cannot receive the packet after K transmission attempts

$$\Pr \{e_{K+1} = Q\} = \sum_{\forall \mathbf{e} \in \{T, Q\}^K} \Pr \{e_{K+1} = Q | \mathbf{e}\} \Pr \{\mathbf{e}\} \quad (1)$$

ANALYTICAL EXPRESSIONS (CONT'D)

- ▶ Applying the chain rule

$$\Pr \left\{ \bigcap_{i=0}^K e_i \right\} = \Pr \{e_0\} \prod_{k=1}^K \Pr \left\{ e_k \mid \bigcap_{i=0}^{k-1} e_i \right\} \quad (2)$$

- ▶ in which

$$\Pr \left\{ e_k = T \mid \bigcap_{i=0}^{k-1} e_i \right\} = 1 - \prod_{\substack{l \in \Phi_T(\mathbf{e}) \\ l < k}} P_{R_l \rightarrow R_k}, \quad (3)$$

$$\Pr \left\{ e_k = Q \mid \bigcap_{i=0}^{k-1} e_i \right\} = \prod_{\substack{l \in \Phi_T(\mathbf{e}) \\ l < k}} P_{R_l \rightarrow R_k}. \quad (4)$$

- ▶ Assuming that $P_{R_u \rightarrow R_l} = 0$ when $l = 0, u < l$

$$\Pr \{ \mathbf{e} \} = \prod_{\substack{\forall l \in \Phi_T(\mathbf{e}) \\ k \in \Theta_Q(\mathbf{e}) \\ k > l}} P_{R_l \rightarrow R_k} \left(1 - \prod_{\substack{u \in \Phi_T(\mathbf{e}) \\ u < l}} P_{R_u \rightarrow R_l} \right) \quad (5)$$

ANALYTICAL EXPRESSIONS (CONT'D)

$$\begin{aligned}
 \Pr \{e_{K+1} = Q\} &= \sum_{\forall \mathbf{e} \in \{T, Q\}^K} \Pr \{e_{K+1} = Q | \mathbf{e}\} \Pr \{\mathbf{e}\} \\
 &= \sum_{\forall \mathbf{e} \in \{T, Q\}^K} \prod_{\substack{\forall l \in \Phi_T(\mathbf{e}) \\ k \in \Theta_Q(\mathbf{e}) \cup \{K+1\} \\ k > l}} P_{R_l \rightarrow R_k} \left(1 - \prod_{\substack{u \in \Phi_T(\mathbf{e}) \\ u < l}} P_{R_u \rightarrow R_l} \right) \quad (6)
 \end{aligned}$$

NUMERICAL RESULTS

Simulation parameters, $N = 12, D = 20m$

Parameter m in Nakagami- m distribution, $PL = 2.32$

Distance between vehicles (meter)	Parameter m
20	1.94
40 to 60	1.86
80 to 160	0.45
180 to 240	0.32

Description	Value
Transmit Power	20 dBm
Antenna gain	4.5 dB
Cable loss	3.4 dB
Noise floor	-99 dBm
SNR decoding threshold	8 dB



IEEE Std 802.11-2012, *IEEE Standard for Information technology - Telecommunications and information exchange between systems - Local and metropolitan area networks - Specific requirements - Part 11: Wireless LAN Medium Access Control (MAC) and Physical Layer (PHY) Specifications*, 2012.



D. Jiang, Q. Chen, and L. Delgrossi, "Optimal data rate selection for vehicle safety communications," in *Proc. ACM Int. Wksp. VANET*, San Francisco, CA, Sep. 2008, pp. 30–38.

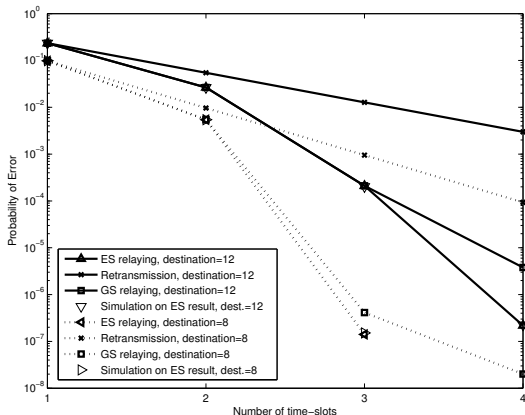


L. Cheng, B. E. Henty, D. D. Stancil, F. Bai, and P. Mudalige, "Mobile vehicle-to-vehicle narrow-band channel measurement and characterization of the 5.9 GHz dedicated short range communication (DSRC) frequency band," *IEEE Sel. Areas Commun.*, vol. 25, pp. 1501–1516, Aug. 2007.



NUMERICAL RESULTS (CONT'D)

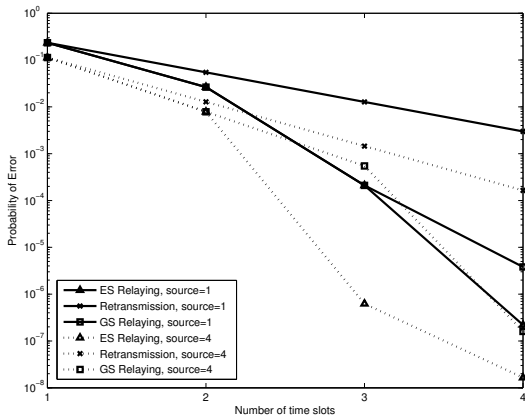
A message that needs to be propagated to a specific receiver



ES = Exhaustive Search, GS = Greedy Search

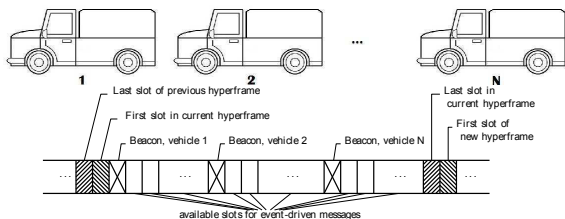
NUMERICAL RESULTS (CONT'D)

A message that originates from one node and needs to be received by the entire platoon, minimising the largest error rate in the platoon



ES = Exhaustive Search, GS = Greedy Search

CONCLUSIONS I



- ▶ Provide a centralized TDMA approach where dissemination of event-driven messages is assisted by a set of relay nodes
- ▶ Give full analysis on the targeted error probability which also includes the best set of relay nodes for a specific dissemination model

CONCLUSIONS II

- ▶ Verified by Monte Carlo simulation using the channel model as suggested from literature.
- ▶ Evaluate the performance of the proposed scheme, it is useful to determine how a message should be spread within a platoon given a strict time-limit
- ▶ Future works
 - ▶ Generalize the probability analysis, not only for distinct relayers
 - ▶ Consider other channel models instead of using a stationary matrix \mathbf{P}

Thank you for your attention